Conversion of Modular Numbers to their Mixed Radix Representation by a Matrix Formula

By J. Schönheim

Introduction. Let $m_i > 1$, $(i = 1, 2, \dots, s)$, be integers relatively prime in pairs and denote $m = m_1 m_2 \cdots m_s$. If x_i , $0 \leq x_i < m_i$, $(i = 1, 2, \dots, s)$ are integers, the ordered set (x_1, x_2, \dots, x_s) is called a modular number, with respect to the moduli m_i $(i = 1, 2, \dots, s)$ and it denotes a unique residue class mod m.

Modular arithmetic has been developed [1], [2], [5], and its use in computers has been suggested [1], [5]. It has also been applied in the solution of various problems [2], [6].

A central question is to determine the least nonnegative residue mod m of a given residue class (x_1, x_2, \dots, x_s) . Denote it by n. In order to work entirely in the given modular system it was suggested [1], [3], [7] and [8] to obtain n in its mixed radix representation with respect precisely to the radices m_i $(i = 1, 2, \dots, s)$, thus in the form

$$n = b_1 + b_2 m_1 + b_3 m_1 m_2 + \cdots + b_s m_1 m_2 \cdots m_{s-1}$$

where $0 \leq b_i < m_i$, $(i = 1, \dots, s)$. In these methods the modular number (b_1, b_2, \dots, b_s) is obtained from the modular number (x_1, x_2, \dots, x_n) sequentially or iteratively.

We propose here (see Theorem) a matrix method which consists in precalculating (s-1) matrices, A_i , $(i = 1, 2, \dots, s-1)$, which depend only on the moduli m_i $(i = 1, 2, \dots, s)$ and in obtaining (b_1, b_2, \dots, b_s) by postmultiplication of (x_1, x_2, \dots, x_s) by A_1, A_2, \dots, A_{s-1} or more precisely, observing the nonassociativity of the used matrix product, computing:

 $(b_1, b_2, b_3, \cdots, b_s) = [\cdots [[(x_1, x_2, x_3, \cdots, x_s)A_1]A_2] \cdots A_{s-2}]A_{s-1}.$

This method is simpler than Mann's method [3] and concentrates the sequential Svoboda-Lindamood-Shapiro method [1], [4] in a single matricial formula.

Definition 1. Let $A = [a_{ij}]$ and $B = [b_{ij}]$ be matrices of s columns with integer elements, whose rows may be regarded as modular numbers with respect to the moduli m_i $(i = 1, \dots, s)$. Define, provided B has s rows, C = AB as $C = [c_{ij}]$, $c_{ij} = \sum a_{ij}b_{ij} \pmod{m_j}, 0 \leq c_{ij} < m_j$.

This matrix multiplication is not associative in general, but two exceptions are mentioned in the following lemma.

LEMMA 1. Let $E = E_{i\nu(c_{\nu})}$ (fixed $i, \nu = 1, 2, \dots, h < s$) be $s \times s$ matrices having units in the main diagonal, c_{ν} as ν th element in the ith ($\nu \neq i$) row and zeroes elsewhere. Let D be a diagonal matrix of the same size. Then if X is an arbitrary matrix with s columns and A an arbitrary $s \times s$ matrix, we have:

$$(1) \qquad (XA)D = X(AD),$$

(2)
$$(\cdots ((XE_1)E_2)\cdots)E_h = X((\cdots ((E_1E_2)E_3)\cdots)E_h).$$

Proof. Properties (1) and (2) are immediate consequences of the definitions. Received May 18, 1966.

J. SCHÖNHEIM

Remark 1. The matrices E_{ν} ($\nu = 1, \dots, h$) are generalized elementary matrices. Notation. Denote $x = (x_1, x_2, \dots, x_s)$ if x is an arbitrary number of the residue class $(x_1, x_2, \dots, x_s) \mod m$ and denote $n \equiv (x_1, x_2, \dots, x_s)$ if n is the least non-negative residue of the class.

LEMMA 2. If (x_1, x_2, \dots, x_s) is a modular number with respect to the moduli m_i $(i = 1, \dots, s)$ and $n \equiv (x_1, x_2, \dots, x_s)$ while

$$\left(\frac{x_2-x_1}{m_1},\frac{x_3-x_1}{m_1},\cdots,\frac{x_s-x_1}{m_1}\right)$$

means a modular number with respect to the moduli m_i ($i = 2, 3, \dots, s$) then

$$\frac{n-x_1}{m_1} \equiv \left(\frac{x_2-x_1}{m_1}, \frac{x_3-x_1}{m_1}, \cdots, \frac{x_s-x_1}{m_1}\right).$$

Proof. $n - x_1$ is divisible by m_1 and since $0 \leq n < m$, it follows that

$$0 \leq \frac{n-x_1}{m_1} < \frac{m}{m_1}.$$

Definition 2. Let $m_i^{-1} \equiv m_{ij} \pmod{m_j}$, $i < j \leq s, 0 < m_{ij} < m_j$ and put $n_{ij} = m_j - m_{ij}$. Let I_k be the identity matrix of rank k. Define, for $1 \leq k \leq s - 1$, $s \times s$ matrices,

	$\int I_{k-1}$	0				
$A_k =$	0	1 0 0	${n_{k,k+1} \atop m_{k,k+1} \over 0}$	$egin{array}{c} n_{k,k+2} \ 0 \ m_{k,k+2} \end{array}$	 	$egin{array}{c} n_{k,s} \ 0 \ 0 \end{array}$
		: 0	0	0	· · •	m_{ks}

LEMMA 3. If (y_1, y_2, \dots, y_s) is a modular number with respect to the moduli m_i $(i = 1, \dots, s)$, then

(3)
$$(y_1, y_2, \cdots, y_s) A_k = \left(y_1, y_2, \cdots, y_k, \frac{y_{k+1} - y_k}{m_k}, \cdots, \frac{y_s - y_k}{m_k}\right).$$

Proof. The matrix A_k is the product of the elementary matrices $E_{k,k+1}(n_{k,k+1})$ $\cdots E_{ks}(n_{ks})$ multiplied by the diagonal matrix



By Lemma 1 associativity holds and the effect of postmultiplication by A_k is the

254

same as the effect of successive postmultiplications by $E_{k,k+1}$, $E_{k,k+2}$, \cdots , E_{ks} and D, which is precisely the right side of (3).

LEMMA 4. Let $n \equiv (x_1, x_2, \dots, x_s)$ and let $q_i, r_i (i = 1, \dots, s)$ be the quotients and the remainders in the successive divisions

(4)
$$n = m_1 q_1 + r_1,$$

$$q_i = m_{i+1}q_{i+1} + r_{i+1}$$
 $(i = 1, \dots, s - 1)$

then

 $(\cdots((x_1, x_2, \cdots, x_s)A_1)A_2)\cdots)A_k = (r_1, r_2, \cdots, r_k, r_{k+1}, y_{k+2}, y_{k+3}, \cdots, y_s)$ and

$$(r_{k+1}, y_{k+2}, \cdots, y_s) \equiv q_k$$
 .

Proof. Proceed by induction on k. Let k = 1. Then by Lemma 3

$$(x_1, \cdots, x_s)A_1 = \left(x_1, \frac{x_2 - x_1}{m_1}, \cdots, \frac{x_2 - x_1}{m_2}\right),$$

hence $r_1 = x_1$ and by Lemma 2,

$$\left(\frac{x_2-x_1}{m_1}, \cdots, \frac{x_s-x_1}{n_1}\right) \equiv \frac{n-x_1}{m_1} = q_1.$$

Therefore

$$\frac{x_2 - x_1}{m_1} \equiv r_2 \pmod{m_2} \quad 0 \le r_2 < m_2.$$

Suppose the assertion is true for $1 < k < h \leq s - 1$, thus

(5) $(\cdots((x_1, x_2, \cdots, x_s)A_1)\cdots)A_{h-1} = (r_1, r_2, \cdots, r_h, y_{h+1}, y_{h+r}, \cdots, y_s),$ and

(6)
$$q_{h-1} \equiv (r_h, y_{h+1}, \cdots, y_s)$$

with respect to the moduli m_i $(i = h, h + 1, \dots, s)$. Then by Lemma 3 and (5)

$$((\cdots ((x_1, x_2, \cdots, x_s)A_1) \cdots)A_{h-1})A_h = \left(r_1, r_2, \cdots, r_h \frac{y_{h+1} - r_h}{m_h}, \cdots, \frac{y_s - r_h}{m_h}\right)$$

and by (6) and Lemma 2

$$\left(\frac{y_{h+1}-r_h}{m_h},\cdots,\frac{y_s-r_h}{m_h}\right)\equiv\frac{q_{h-1}-r_h}{m_h}=q_h.$$

Therefore

$$\frac{y_{h+1} - r_h}{m_h} = r_{h+1}, \qquad 0 \le r_{h+1} < m_{h+1}$$

Hence the result is true for k = h.

THEOREM. If $m_i, m_i > 1$ $(i = 1, 2, \dots, s)$ are integers, relatively prime in pairs

J. SCHÖNHEIM

 $m = m_1 \cdots m_s$, and if n is the least nonnegative residue of the class $(x_1, x_2, \cdots, x_s) \mod m$ and b_1, b_2, \cdots, b_s are the digits of the mixed radix representation of n with respect to the radices m_i $(i = 1, \cdots, s)$ then with matrix multiplication and matrices A_i $(i = 1, \cdots, s)$ as defined in Definitions 1 and 2

$$(b_1, b_2, \cdots, b_s) = (\cdots (((x_1, x_2, \cdots, x_s)A_1)A_2)\cdots)A_{s-1}$$

Proof. The digits b_1, \dots, b_s of the required representation are the remainders of the successive divisions (4) and the theorem is a corollary of Lemma 4 with k = s - 1.

Remark 2. The above algorithm requires in general s - 1 matrix multiplications, but if k < s - 1 and

(7)
$$(\cdots (((x_1, x_2, \cdots, x_s)A_1)A_2)\cdots)A_k = (r_1, r_2, \cdots, r_{k+1}, 0, 0, \cdots, 0)$$

then the right side of (7) is the result, and no further multiplications are needed.

Example. Let 2, 3, 5, 7 be the moduli m_1 , m_2 , m_3 , m_4 . Then the numbers m_{ij} , i < j are given by

and therefore the numbers n_{ij} are

$$\begin{array}{cccc}1&2&3\\&3&2\\&&4\end{array}$$

The matrices A_1 , A_2 , A_3 are

$$A_{1} = \begin{bmatrix} 1 & 1 & 2 & 3 \\ 0 & 2 & 0 & 0 \\ 0 & 0 & 3 & 0 \\ 0 & 0 & 0 & 4 \end{bmatrix}; \quad A_{2} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 3 & 2 \\ 0 & 0 & 2 & 0 \\ 0 & 0 & 0 & 5 \end{bmatrix}; \quad A_{3} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 4 \\ 0 & 0 & 0 & 3 \end{bmatrix}$$

Let (0200) be a residue class mod 210. Let *n* be the least nonnegative residue of this class. Then b_1 , b_2 , b_3 , b_4 , the digits of the mixed radix representation of *n*, with respect to the radices 2, 3, 5, 7 are given by

$$(b_1, b_2, b_3, b_4) = (((0\ 2\ 0\ 0)A_1)A_2)A_3 = (0\ 1\ 3\ 4).$$

Indeed $0 + 1 \cdot 2 + 3 \cdot 2 \cdot 3 + 4 \cdot 2 \cdot 3 \cdot 5 = 140$, 140 < 210 and $140 \equiv 0 \pmod{2}$, $2 \pmod{3}$, $0 \pmod{5}$ and $0 \pmod{7}$.

Department of Applied Mathematics Tel-Aviv University Ramat-Aviv, Tel-Aviv Israel

256

1. M. VALACH & A. SVOBODA, "Circuit operators," Stroje na. Zpracovani Informaci Sb.,

v. 111, 1957, pp. 247-297. (Czech)
2. H. S. SHAPIRO, "Some remarks on modular arithmetic and parallel computation," *Math. Comp.*, v. 16, 1962, pp. 218-222. MR 26 #4511.
3. H. B. MANN, "On modular computation," *Math. Comp.*, v. 15, 1961, pp. 190-192. MR 22

*** 10944**.

* 10944.
4. G. E. LINDAMOOD & G. SHAPIRO, "Magnitude comparison and overflow detection in modular arithmetic computers," SIAM Rev., v. 5, 1963, pp. 342-350. MR 29 * 6662.
5. A. SVOBODA, "The numerical system of residual classes in mathematical machines," Information Processing, pp. 419-422, UNESCO, Paris, R. Oldenbourg, Munich and Butterworths, London, 1960. MR 28 * 5581.
6. J. BOROSH & A. S. FRAENKEL, "Exact solutions of linear equations with rational coefficients by congruence techniques," Math. Comp., v. 20, 1966, pp. 107-112.
7. V. N. TEITEL'BAUM, "Comparison of numbers in the Czech system of numbers," Dokl. Akad. Nauk SSSR, v. 121, 1958, pp. 807-810. (Russian) MR 21 * 3367.
8. A. S. FRAENKEL, On Size of Modular Numbers, Proc. ACM 19th National Conference, Philadelphia, Pa., 1964.

Philadelphia, Pa., 1964.